

Topological Universality of the Art Gallery Problem

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Abstract

We prove that any compact semi-algebraic set is homeomorphic to the solution space of some art gallery problem. Previous works have established similar universality theorems, but holding only up to homotopy equivalence, rather than homeomorphism, and prior to this work, the existence of art galleries even for simple spaces such as the Möbius strip or the three-holed torus were unknown. Our construction relies on an elegant and versatile gadget to copy guard positions with minimal overhead, and is thus simpler than previous constructions, consisting of a single rectangular room with convex slits cut out from the edges. We additionally show that both the orientable and non-orientable surfaces of genus n require galleries with only $O(n)$ vertices.

1 Introduction

An instance of the *art gallery problem* consists of a polygon P (which we refer to as the *art gallery*), and the objective is to find a finite set of points $G \subseteq P$ (the *guards*) of minimal cardinality such that every point in P is *visible* to some guard, meaning that the line segment between the point and the guard is contained within P . This problem was first introduced by Viktor Klee in 1973 and has a long history. In 1986, Lee and Lin [6] showed that the decision problem of determining whether there exists a configuration with at most k guards is NP-hard, but the problem is not known to be in NP since there is no obvious succinct way to represent the guards' coordinates in a solution (they might have to be irrational, even if all polygon vertex coordinates are rational). In 2018, Abrahamsen, Adamaszek, and Miltzow [1] showed that the problem is $\exists\mathbb{R}$ -complete, mostly settling the question of complexity (modulo the longstanding conjectures that $\text{NP} \subsetneq \exists\mathbb{R} \subsetneq \text{PSPACE}$). Approximation algorithms and lower bounds have also been studied; see Vaezi and Ghodsi [10] for a recent overview.

Aside from complexity aspects, a parallel line of inquiry concerns the topology of the space of solutions. Supposing that gallery P requires exactly k guards at minimum, we let

$$V(P) := \{G \subseteq P \mid |G| = k \text{ and every } p \in P \text{ is visible to some } g \in G\}.$$

The set $V(P)$ can consist of unordered sets of points of cardinality k , which can be turned into a natural metric space using the Hausdorff distance

$$d_H(G_0, G_1) := \max_{i \in \{0,1\}} \max_{g \in G_i} \min_{g' \in G_{1-i}} d(g, g'),$$

where d denotes the Euclidean distance on \mathbb{R}^2 . Thus, for any art gallery P , $V(P)$ is a topological space, so a natural question is, what kinds of topological spaces can occur?

Aside from being a mathematical curiosity, this question is relevant to complexity theory due to the connections between $\exists\mathbb{R}$ -hardness and topological universality. The most famous example is Mnëv’s theorem [8] that any semi-algebraic set is homotopy-equivalent to the space of point configurations of some order type, inspiring an eventual proof that order type realizability is $\exists\mathbb{R}$ -complete [9]. A *semi-algebraic* set is a finite union $\bigcup_{i=1}^m S_i$, where each S_i is defined to be the set of points in $\mathbb{R}^n (x_1, x_2, \dots, x_n)$ satisfying a finite number of constraints of form $P(x_1, x_2, \dots, x_n) \geq 0$ or $P(x_1, x_2, \dots, x_n) > 0$, where P is a polynomial. The canonical complete problem for $\exists\mathbb{R}$ is called ETR (Existential Theory of the Reals), which asks whether a given semi-algebraic set is nonempty. As a consequence of their reduction from ETR to the art gallery problem, Abrahamsen et al. [1] show, for any compact semi-algebraic set S , how to construct an art gallery P such that $V(P)$ surjects continuously onto S . However, they do not show that the mapping is injective, so this fails to establish universality. In a recent paper, Bertschinger, El Maalouly, Miltzow, Schnider, and Weber [2] show that any compact semi-algebraic set S is homotopy-equivalent to $V(P)$ for some polygon P . They leave as an open question whether P can be constructed so that $V(P)$ is not just homotopy-equivalent, but homeomorphic to S , with only the following list of spaces shown to be captured up to homeomorphism:

- k -clover (obtained by joining k circles at a single point)
- k -chain (obtained by connecting k circles in a path with $k - 1$ disjoint line segments)
- $4k$ -necklace (obtained by connecting k circles in a cycle with k disjoint line segments)
- k -sphere
- The torus
- The 2-holed torus

The constructions for these spaces are all based on simple galleries with solution spaces homeomorphic to a circle and an interval, which are combined to give Cartesian products and then given simple additional constraints. However, using these methods it is difficult to obtain more general spaces because the geometry significantly limits the constraints that can be used. Thus, prior to this work, homeomorphism universality was unknown even for closed surfaces. Galleries for the real projective plane, Klein bottle, and Möbius strip were explicitly left as open questions.

In this work, we settle the question of homeomorphism universality in the affirmative:

Theorem 1.1. *For every compact semi-algebraic set $S \neq \emptyset$, there exists a polygon P such that S is homeomorphic to $V(P)$.*

In addition to yielding a strictly stronger result than that of Bertschinger et al. [2], our construction is structurally simpler: the art gallery always consists of a single rectangular room with two types of convex slits repeatedly cut out of the edges. As an example, we explicitly draw an art gallery for the Möbius strip (to the best resolution that can reasonably fit on a page) in Figure 5.

The key ingredient in our approach is a novel gadget which “copies” a variable (in effect adding a constraint of form $x_i = x_j$). Gadgets with similar functions can also be found in Abrahamsen et al. [1]. Our gadget is an improvement for the following reasons:

- It is more versatile. In particular the polygon P does not need to intersect the convex hull of the segments being copied.

- It works in more general contexts. The previous gadget uses constraints created by the need for guards to see every point on the interior of the polygon P . Our gadget works in the variant of the problem where guards only need to see the boundary of the polygon.

Our general construction, which we present in Section 2, does not yield an obvious bound on the size of P (number of vertices) as a function of the complexity of the space S . In Section 3, we refine our construction to show that both the orientable and non-orientable surfaces of genus n can be captured by art galleries with only $O(n)$ vertices.

2 General construction

As in Bertschinger et al. [2], the starting point for our reduction is a theorem by Hironaka [4] that any compact semialgebraic set can be triangulated as a cubical complex. We begin by reviewing this result. Next, we introduce our variable copying gadget and establish its key properties. Finally, we show how to combine the gadgets to construct the art gallery. We illustrate our construction using the Möbius strip as a running example.

2.1 Hironaka’s theorem

An *abstract cubical complex* is subset K of a set

$$I^n = \left\{ \{\mathbf{x} \in [0, 1]^n \mid x_{i_1} = c_1, \dots, x_{i_k} = c_k\} \mid 1 \leq i_1 < \dots < i_k \leq n, (c_1, \dots, c_k) \in \{0, 1\}^k \right\}$$

of faces of an n dimensional hypercube $[0, 1]^n$ such that if $a \in K$ and $b \in I^n$ and $b \subseteq a$, then $b \in K$. Every cubical complex corresponds to a subset $|K| = \bigcup K$ of $[0, 1]^n$, called its *geometric realization*.

Theorem 2.1 (Bertschinger et al. [2], Lemma 3). *Any compact semi-algebraic set is homeomorphic to the geometric realization of an abstract cubical complex.*

Proof sketch. Using Hironaka’s theorem [4], it is possible to show that any compact semi-algebraic set is homeomorphic to the geometric realization of an abstract simplicial complex (see Hoffman [5]). Additionally, it can be shown that any abstract simplicial complex has an abstract cubic complex with a homeomorphic geometric realization (see Blass and Wlodzimierz [3, Theorem 1.1]). \square

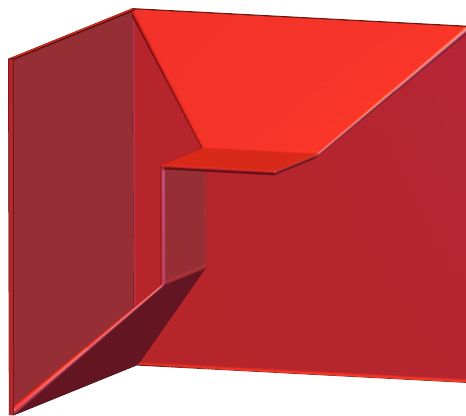


Figure 1: A 4-dimensional geometric realization of a cubical complex homeomorphic to a Möbius strip.

For example, the circle S^1 can be represented as the abstract cubical complex

$$\{\{\mathbf{x} \in [0, 1]^2 \mid x_1 = 0\}, \{\mathbf{x} \in [0, 1]^2 \mid x_1 = 1\}, \{\mathbf{x} \in [0, 1]^2 \mid x_2 = 0\}, \{\mathbf{x} \in [0, 1]^2 \mid x_2 = 1\}\},$$

in which case its geometric realization is the boundary of the unit square in \mathbb{R}^2 . The Möbius strip can be realized as a subset of the 4-dimensional hypercube as in Figure 1.

2.2 The variable and copying gadgets

The art gallery we construct will be an axis-aligned rectangle with several slits removed from around the border. We refer to certain slits and combinations of slits as *gadgets*. Our construction involves three kinds of gadgets: *variable gadgets*, *copying gadgets*, and *clause gadgets*. Our copying gadgets are based on the following geometric observation.

Lemma 2.2. *Let L_1 and L_2 be two distinct parallel lines in the plane, and let Z be a point strictly between L_1 and L_2 . Let $f_Z : L_1 \rightarrow L_2$ be the map defined via inversion through Z , i.e., for $p \in L_1$, $f_Z(p)$ is the intersection of the line pZ with L_2 . Then f_Z is linear (in the sense that it preserves ratios of distances between points on L_1 and the corresponding points on L_2).*

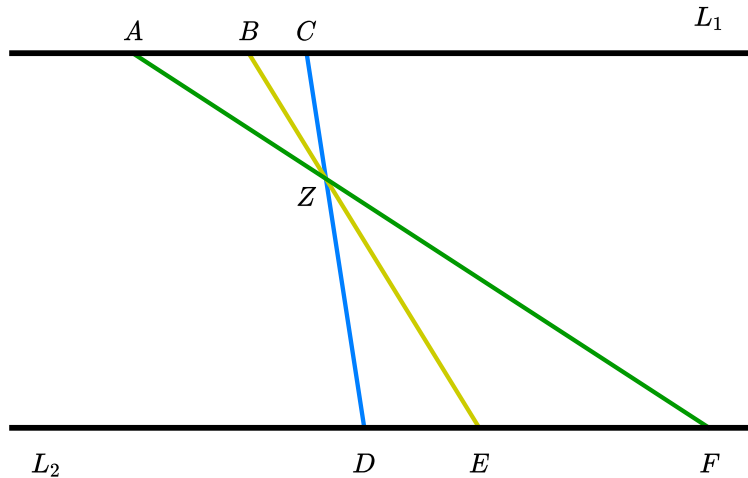


Figure 2

Proof. Without loss of generality, assume L_1 and L_2 are drawn aligned with the x -axis as in Figure 2. Let A , B , and C be three arbitrary distinct points on L_1 , and let F , E , and D be their respective images under f_Z . It is easy to see that (AZB, FZE) and (AZC, FZD) are pairs of similar triangles. This means that the ratios of lengths of AB to FE and AC to FD are equal, that is $\frac{B_x - A_x}{F_x - E_x} = \frac{C_x - A_x}{F_x - D_x}$. Rearranging, we have $\frac{B_x - A_x}{C_x - A_x} = \frac{F_x - E_x}{F_x - D_x}$. \square

The following lemma summarizes the key properties of the variable and copying gadgets, which are depicted in Figure 3.

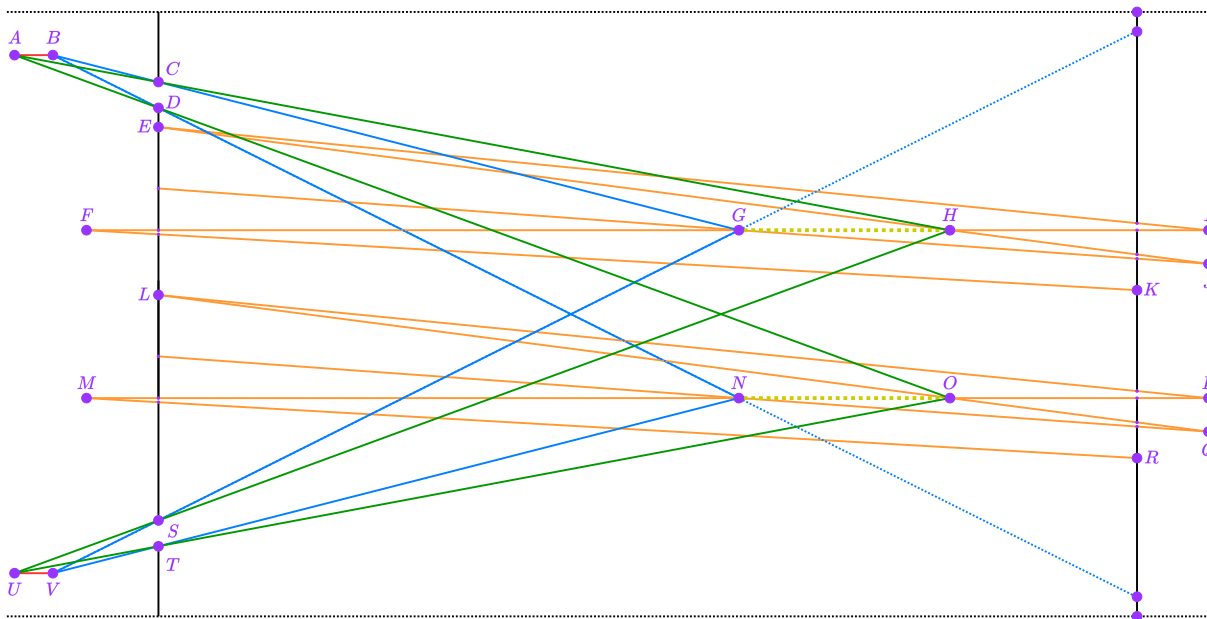


Figure 3: The variable and copying gadgets. The boundary of this section of the art gallery consists of the entire outer profile (continuing above and below the dashed lines at the top and bottom), involving line segments of all colors. A pair of guards are constrained to patrol the line segments GH and NO at the same x -coordinate. There are two variable gadgets, shown in orange. The first one consists of the slits at points F , I , and J , which enforces that one guard must patrol GH . The second consists of the slits at points M , P , and Q , which enforces that the other guard must patrol NO . The other two slits, $CBAD$ and $SUVT$, comprise the copying gadget, which enforces that both guards have the same x -coordinate.

Lemma 2.3. *As long as all co-linearity relations and ordering relations between x and y coordinates of highlighted points are as depicted in Figure 3, with no other slits between the black dashed lines at the top and bottom, it follows that:*

- (1) *Any valid guard placement requires at least two guards between the dashed lines.*
- (2) *Such a configuration of two guards is valid if and only if the guards are placed on line segments GH and NO at the same x -coordinate.*

Proof. To see point F , we must have at least one guard placed within the triangle FIK ; to see point I , we must have at least one guard placed within the triangle EFI ; to see point M , we must have at least one guard placed within the triangle MPR ; and to see point P , we must have at least one guard placed within the triangle LMP . Clearly, there is no single guard that can see all four of these points (as, for instance, triangles EFI and LMP are disjoint), which proves (1). Furthermore, any solution with only two guards must have one guard on line segment FI and another guard on line segment MP , for otherwise at least one of the four triangles will not have a guard in it. Given that the guards are on these line segments, the only way that points J and Q will be visible is if one guard is on segment GH and the other guard is on NO .

Thus, to all that remains to prove (2) is that the two guards must have the same x -coordinate. Let α be the location of the guard on GH , and let β be the location of guard on NO ; our objective is to show that, in any valid configuration, $\alpha_x = \beta_x$. We define the following four additional points:

- $\bar{\alpha} :=$ the intersection of line segment AB with line αC
- $\underline{\alpha} :=$ the intersection of line segment UV with line αS
- $\bar{\beta} :=$ the intersection of line segment AB with line βD
- $\underline{\beta} :=$ the intersection of line segment UV with line βT

Due to the obstructions by points C and D , the guard at α can only see AB to the left of $\bar{\alpha}$, and the guard at β can only see AB to the right of $\bar{\beta}$. If $\bar{\alpha}_x < \bar{\beta}_x$, the line segment $\bar{\alpha}\bar{\beta}$ will therefore not be seen by any guard, so we must have that $\bar{\alpha}_x \geq \bar{\beta}_x$. Similarly, the guard at α can only see UV to the right of $\underline{\alpha}$, and the guard at β can only see UV to the left of $\underline{\beta}$, so we must have that $\underline{\alpha}_x \leq \underline{\beta}_x$. Thus, in any valid configuration of the two guards,

$$\begin{aligned}
\frac{\alpha_x - G_x}{H_x - G_x} &= \frac{B_x - \bar{\alpha}_x}{B_x - A_x} && \text{(by Lemma 2.2 with } Z := C\text{)} \\
&\leq \frac{B_x - \bar{\beta}_x}{B_x - A_x} \\
&= \frac{\beta_x - N_x}{O_x - N_x} && \text{(by Lemma 2.2 with } Z := D\text{)} \\
&= \frac{V_x - \underline{\beta}_x}{V_x - U_x} && \text{(by Lemma 2.2 with } Z := T\text{)} \\
&\leq \frac{V_x - \underline{\alpha}_x}{V_x - U_x} \\
&= \frac{\alpha_x - G_x}{H_x - G_x} && \text{(by Lemma 2.2 with } Z := S\text{)}.
\end{aligned}$$

Multiplying through by $(H_x - G_x)$ and using the fact that $N_x = G_x$ and $O_x = H_x$, we have

$$\alpha_x - G_x \leq \beta_x - G_x \leq \alpha_x - G_x.$$

Thus, $\alpha_x = \beta_x$ in any valid configuration (and it is easy to check that such a configuration is indeed valid), which proves (2). \square

2.3 Constructing the art gallery

We are now ready to prove homeomorphism universality.

Proof of Theorem 1.1. By Theorem 2.1, we may assume without loss of generality that S is the geometric realization of a nonempty cubical complex. This means that S can be described as a subset of $[0, 1]^n$ whose coordinates x_1, x_2, \dots, x_n satisfy some formula

$$\phi_S \equiv \bigvee_{j=1}^{m'} ((x_{i'_j,1} = c'_{j,1}) \wedge (x_{i'_j,2} = c'_{j,2}) \wedge \dots \wedge (x_{i'_j,\ell'} = c'_{j,\ell'})),$$

where each $i'_{j,k} \in \{1, 2, \dots, n\}$ and $c'_{j,k} \in \{0, 1\}$. The formula will always be in disjunctive normal form (DNF), where each conjunctive clause corresponds to a face of the cubical complex, so ϕ_S describes the set of points lying in at least one such face. For example, the formula for the Möbius strip M depicted in Figure 1 has six clauses, one for each of the six 2-dimensional faces. To write an explicit formula ϕ_M , we choose a coordinate system where $x_4 = 0$ corresponds to the points on

the outer shell and enumerate the faces as clauses by traversing the strip starting from the left-most depicted face and proceeding next toward the back face:

$$\begin{aligned}\phi_M \equiv & (x_2 = 0 \wedge x_4 = 0) \vee (x_1 = 0 \wedge x_4 = 0) \vee (x_1 = 0 \wedge x_3 = 1) \\ & \vee (x_3 = 1 \wedge x_4 = 1) \vee (x_2 = 0 \wedge x_4 = 1) \vee (x_2 = 0 \wedge x_3 = 0)\end{aligned}$$

We next rewrite ϕ_S in conjunctive normal form (CNF),

$$\phi_S \equiv \bigwedge_{j=1}^m ((x_{i_{j,1}} = c_{j,1}) \vee (x_{i_{j,2}} = c_{j,2}) \vee \cdots \vee (x_{i_{j,\ell}} = c_{j,\ell})),$$

by enumerating all tuples of constraints that take one constraint from each clause. For example, since ϕ_M has 6 clauses in DNF, each of size 2, the standard translation to CNF produces $2^6 = 64$ clauses, each of size 6. However, after eliminating redundancies, we can simplify ϕ_M to

$$\begin{aligned}\phi_M \equiv & (x_1 = 0 \vee x_2 = 0 \vee x_3 = 1) \\ & \wedge (x_2 = 0 \vee x_3 = 1 \vee x_4 = 0) \\ & \wedge (x_1 = 0 \vee x_2 = 0 \vee x_4 = 1) \\ & \wedge (x_3 = 0 \vee x_3 = 1 \vee x_4 = 0 \vee x_4 = 1) \\ & \wedge (x_1 = 0 \vee x_3 = 0 \vee x_4 = 0 \vee x_4 = 1).\end{aligned}\tag{1}$$

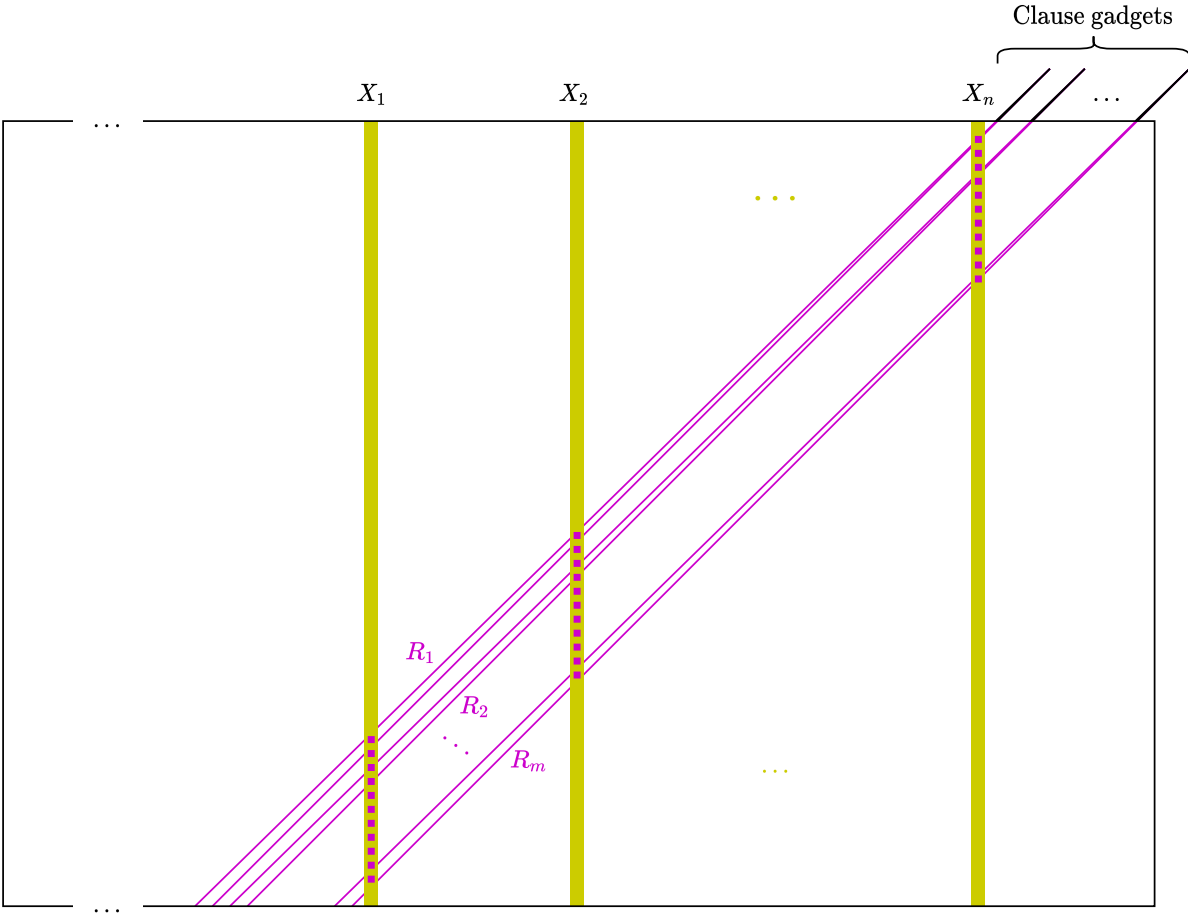


Figure 4: Illustration of the first step of the general construction (making clause gadgets), starting from a CNF formula with n variables and m clauses. The full art gallery is shown in more detail for the Möbius strip in Figure 5.

Starting from a rectangular art gallery, we make a narrow diagonal slit for each clause j in the top-right corner such that the regions R_j of the gallery that can see to the end of each slit extend downward to the left and do not overlap, as shown in Figure 4. For each $i \in \{1, 2, 3, \dots, n\}$ we define X_i be tall, skinny, axis-aligned rectangular region such that the convex hulls of the sets $X_i \cap (R_1 \cup R_2 \cup \dots \cup R_m)$ do not overlap in y -coordinates (in terms of Figure 4, the pink dashed lines must not overlap in y -coordinates). It is always possible to guarantee this non-overlapping property by making the gallery sufficiently tall and/or making the clauses sufficiently close together.

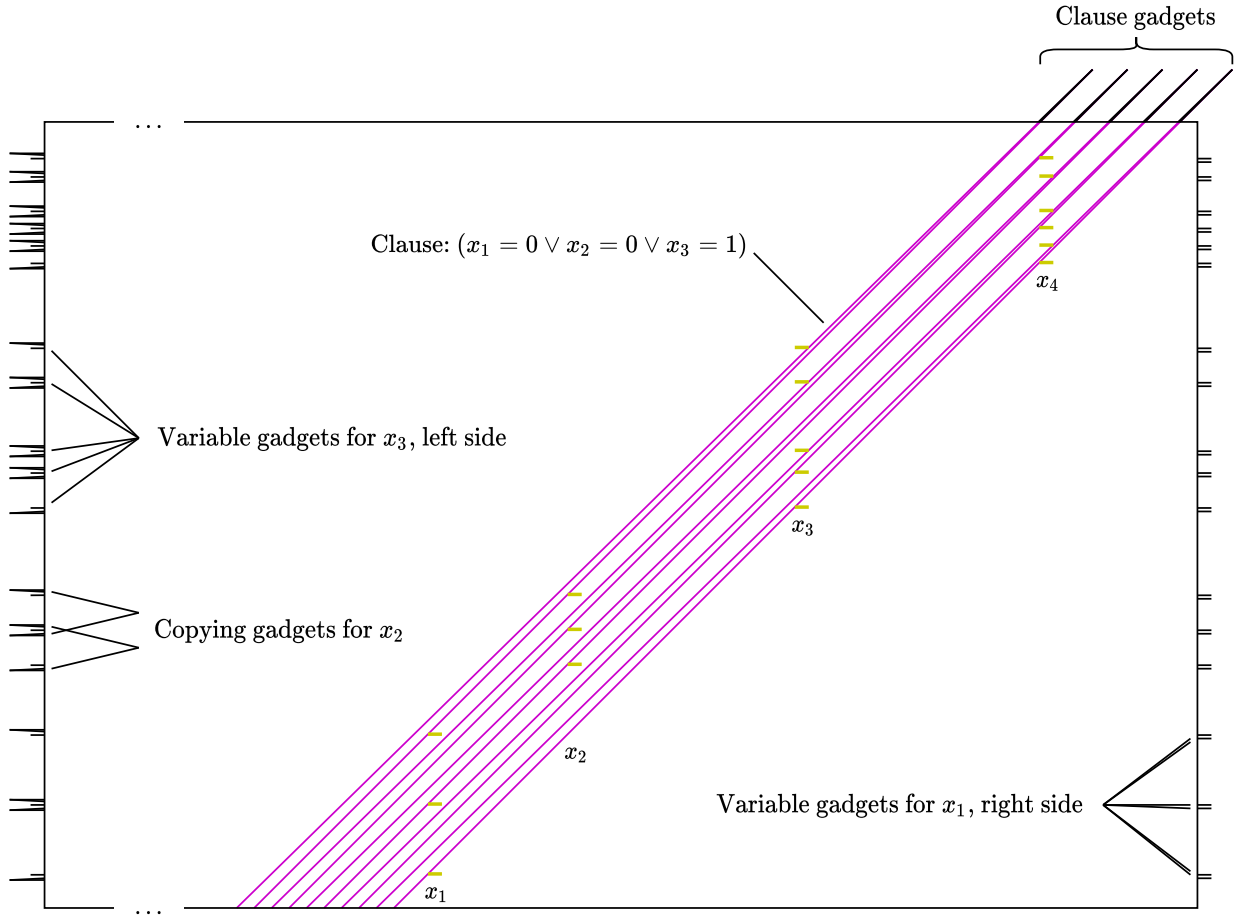


Figure 5: An art gallery whose solution space is homeomorphic to a Möbius strip, using the formula ϕ_M from (1). The horizontal gold dashes above each variable are the line segments on which guards must walk, as in Figure 3. The pink diagonal lines depict the regions R_j that can see to the end of each clause gadget slit; each of these regions must contain at least one guard. The left wall must be placed sufficiently far to the left to ensure that none of the variable/copy gadgets interfere with each other. As drawn, this particular art gallery has 183 vertices.

In every clause j , for every constraint $x_{i_j,k} = c_{j,k}$, we add a variable gadget in the left and right sides of the rectangle enforcing the constraint that there is a guard on a short horizontal line segment (hereafter a *guard segment*) spanning the width of $X_{i_j,k}$ whose endpoint lies within R_j . If $c_{j,k} = 0$, this will be the left endpoint, and if $c_{j,k} = 1$, this will be the right endpoint. We then add copying gadgets to enforce that all guard segments placed in X_i for the same i must have the same x -coordinate. If the left wall is moved sufficiently far to the left and the cuts are made sufficiently small, we can realize the hypothesis of Lemma 2.3 that the horizontal slice including every pair of consecutive copying gadgets is drawn as in Figure 3, without any other guard segments appearing between the dashed lines. It then follows from repeatedly applying Lemma 2.3 that, letting P_0 be the polygon P without the clause gadget slits, $V(P_0)$ consists of all solutions with one guard on each guard segment, with guards within the same X_i placed at the same x -coordinate. There is thus a natural homeomorphism $h : V(P_0) \rightarrow [0, 1]^n$, and from the way the clause gadgets were

constructed, it clearly follows that $V(P)$ consists of all solutions in which ϕ_S is satisfied under h . Thus, $V(P)$ is homeomorphic to S . \square

3 Efficient construction for closed surfaces

Even if a space can be triangulated as a cubical complex with relatively few faces, the conversion from DNF to CNF can exponentially blow up the number of clauses, and thus the size of the art gallery. Here we show that an important class of topological spaces, namely the closed surfaces, can occur as solutions to art galleries with linearly many vertices.

Theorem 3.1. *There are polygons P_n, Q_n with $O(n)$ vertices such that $V(P_n)$ is homeomorphic to the closed orientable surface of genus n and $V(Q_n)$ is homeomorphic to the closed non-orientable surface of genus n .*

Proof. We know by Theorem 1.1 that such polygons exist for $n \in \{0, 1\}$, but the number of vertices of these galleries have no effect on the asymptotic properties of P_n and Q_n , so we may ignore these cases.

It is well known that, for $n \geq 2$, the orientable surface of genus n can be obtained as the connected sum of n copies of a torus, $T^2 \# T^2 \# \dots \# T^2$, while the non-orientable surface of genus n can be obtained as the connected sum of n copies of the real projective plane $\mathbb{R}P^2 \# \mathbb{R}P^2 \# \dots \# \mathbb{R}P^2$ (see, e.g., the textbook by Massey [7]). The connect sum $R \# R$ is ordinarily defined as removing an open disk from two copies of R and gluing their boundaries together. For 2-dimensional surfaces, this is equivalent to gluing the disk boundaries to opposite ends of a cylindrical tube, which is the formulation we use in this construction. Thus, let R be either T^2 or $\mathbb{R}P^2$. By Theorem 2.1, we know there is some j such that R is homeomorphic to $|C|$, the geometric realization of a cubical complex C with j variables. Let x_1, x_2, \dots, x_j be these variables and write $\mathbf{x} = (x_1, x_2, \dots, x_n)$.

Note that C has at least two 2-dimensional faces. Let C_1 be a cubical complex obtained by removing a face f_1 of dimension 2 from C , and C_2 obtained by removing a different face. For $i \in \{1, 2\}$, define B_i to be the cubical complex containing f_i and its boundary, so that $|C_i| \cap |B_i|$ is the boundary of the removed disc (see Figure 6).

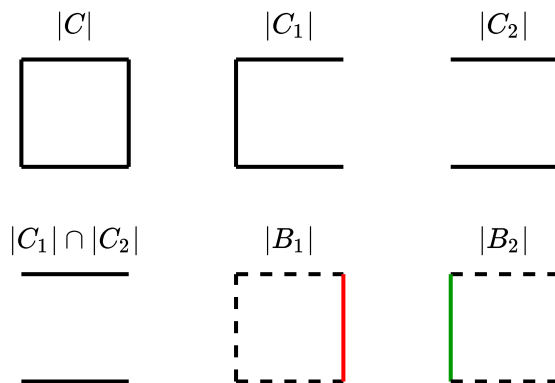


Figure 6: Visualisation of $|C_1|$, $|C_2|$, $|B_1|$ and $|B_2|$. In reality, C should be a complex for T^2 or $\mathbb{R}P^2$, rather than S^1 as shown. The figure is just meant to convey the construction at a schematic level.

Define a new variable x_0 and fix constants $k_0 = 0 < k_1 < \dots < k_{n-2} < 1 = k_{n-1}$. Then the following formulas express the property that a point \mathbf{x} is contained in $R \# R \# \dots \# R$ (see Figure 7

for a visualization). For n even we have

$$\begin{aligned} & (\mathbf{x} \in |C_2| \vee x_0 = 0) \wedge (\mathbf{x} \in |C_1| \vee x_0 = 1) \wedge \\ & (\mathbf{x} \in |B_1| \vee 0 \leq x_0 \leq k_1 \vee k_2 \leq x_0 \leq k_3 \vee \cdots \vee k_{n-3} \leq x_0 \leq k_{n-2} \vee x_0 = 1) \wedge \\ & (\mathbf{x} \in |B_2| \vee k_1 \leq x_0 \leq k_2 \vee k_3 \leq x_0 \leq k_4 \vee \cdots \vee k_{n-2} \leq x_0 \leq 1), \end{aligned}$$

and for n odd we have

$$\begin{aligned} & (\mathbf{x} \in |C_2| \vee x_0 = 0 \vee x_0 = 1) \wedge \mathbf{x} \in |C_1| \wedge \\ & (\mathbf{x} \in |B_1| \vee 0 \leq x_0 \leq k_1 \vee k_2 \leq x_0 \leq k_3 \vee \cdots \vee k_{n-2} \leq x_0 \leq 1) \wedge \\ & (\mathbf{x} \in |B_2| \vee k_1 \leq x_0 \leq k_2 \vee k_3 \leq x_0 \leq k_4 \vee \cdots \vee k_{n-3} \leq x_0 \leq k_{n-2} \vee x_0 = 1). \end{aligned}$$

Note that these spaces are not in general the geometric realization of a cubical complex in $x_0, x_1, x_2, \dots, x_j$ because of the constraints on x_0 ; nevertheless, we will construct art galleries for them.

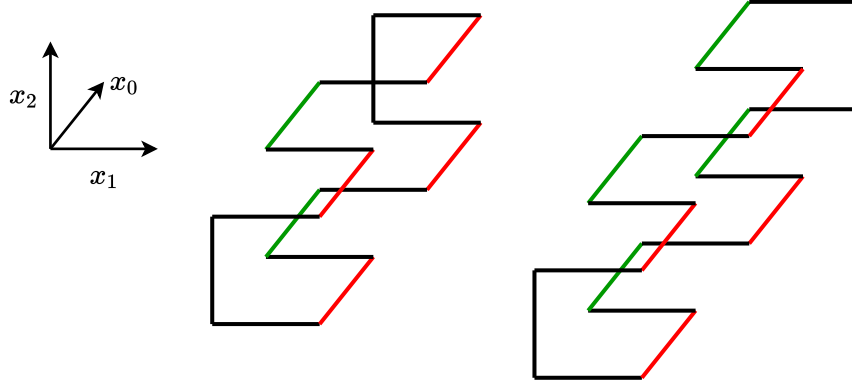


Figure 7: A schematic visualization for our constructions of $R\#R\#\dots\#R$ for even n (left) and odd n (right). For all values of x_0 , we require $\mathbf{x} \in |C_1|$ and $\mathbf{x} \in |C_2|$, except for $x_0 = 0$ or $x_0 = 1$ where one of these constraints is dropped to cap the hole at the end. For $k_i < x_0 < k_{i+1}$, \mathbf{x} must be in either $|B_1| \cap |C_1|$ or $|B_2| \cap |C_2|$ depending on the parity of i . This creates tubes connecting the copies of R .

We can write these expressions in CNF with terms of form $x_i = 0$, $x_i = 1$, or $k_i \leq x_0 \leq k_{i+1}$. We write the constraint $\mathbf{x} \in |B_1|$ in CNF as

$$\phi_{B_1} := \bigwedge_{\ell=1}^p \bigvee_{m=1}^{q_\ell} t_{\ell,m},$$

where each $t_{\ell,m}$ is a term of form e.g. $x_1 = 0$. We may thus rewrite

$$\begin{aligned} & \mathbf{x} \in |B_1| \vee 0 \leq x_0 \leq k_1 \vee k_2 \leq x_0 \leq k_3 \vee \cdots \vee k_{n-2} \leq x_0 \leq 1 \\ & \equiv \bigwedge_{\ell=1}^p \left(\left(\bigvee_{m=1}^{q_\ell} t_{\ell,m} \right) \vee 0 \leq x_0 \leq k_1 \vee k_2 \leq x_0 \leq k_3 \vee \cdots \vee k_{n-2} \leq x_0 \leq 1 \right). \end{aligned}$$

To add the constraints on x_0 to each clause involving x_0 , we must make a small modification to the construction from Section 2. Instead of having the clause gadgets along the top wall, we extend the top wall upward and put these gadgets along the right wall. It is easy to see that this

does not affect our ability to create the constraints on \mathbf{x} . We then add the constraints on x_0 to each clause as in Figure 8; if all clause gadgets are vertically translated copies of each other, then all guard segments for x_0 across the different clauses can be placed at the same x -coordinates.

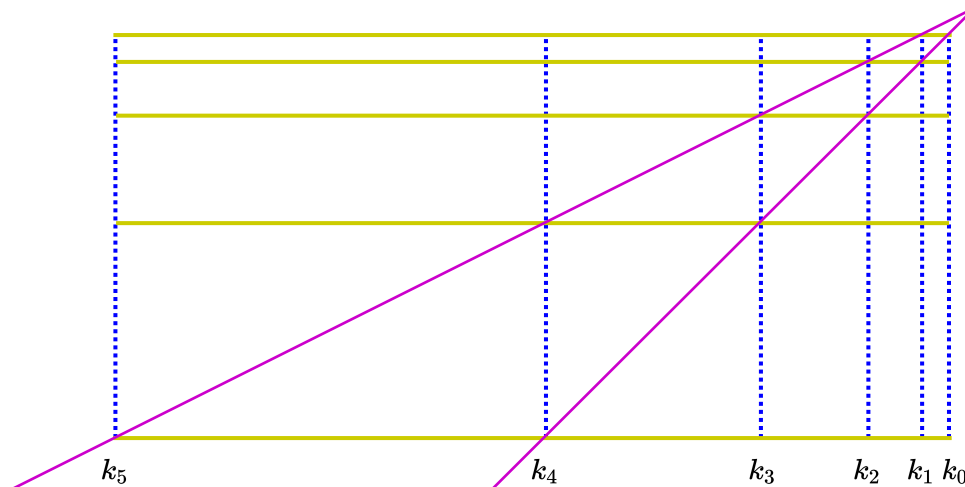


Figure 8: Guard segments for x_0 . We place the first guard segment so that the lower line in the wedge intersects it at a 0, and the upper line intersects it at some point k_1 . We can then place another segment which intersects the *lower* line at k_1 , and the intersection of this segment with the upper line gives the value of k_2 . We can repeat until we have $n - 1$ segments. Each clause depending on x_0 will have a copy of this setup but with some of the segments removed depending on which terms $k_i \leq x_0 \leq k_{i+1}$ it contains.

Using a similar expansion for the other terms, we obtain a CNF expression that cuts out the space we want. Since p and each q_ℓ are constants, the total number of clauses does not depend on n , and each clause has at most $O(n)$ terms. Thus, when we add the constraints involving x_0 to the constant-sized formula ϕ_{B_1} , the number of vertices increases linearly in n .

We may similarly add x_0 constraints to the formula for B_2 with linear blowup. The result is a gallery with $O(n)$ vertices whose solution space is homeomorphic to $R\#R\#\dots\#R$. \square

We have shown that such P_n, Q_n exist having $O(n)$ vertices. In case it is of interest, we leave it to the reader to verify that the positions of the vertices can additionally be chosen to be rational numbers which require only $O(n)$ bits to describe.¹

4 Conclusion

In this work we have settled the open question of Bertschinger et al. [2] by showing that solution spaces to the art gallery problem can capture the topology of any semi-algebraic set up to homeomorphism. In doing so, we have introduced a new form of copying gadget that enables simpler arguments about the structure of valid solutions to art gallery problem instances.

We are hopeful that this gadget may find other applications as well. For example, copying gadgets play a central role in the proof that the art gallery problem is $\exists\mathbb{R}$ -hard [1], and before now

¹In general, if we fix the length of the guard segments for x_0 then the wedge parameters that would give $k_{n-1} = 1$ are not rational numbers. Instead, we fix a sufficiently thin wedge and choose the length of the guard segments for x_0 appropriately.

it was unknown if they could still be created in the setting where guards need only see the walls of the art gallery. Our gadgets could be a useful step towards showing that the wall-only variant is also $\exists\mathbb{R}$ -hard.

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